

# ROLE ACTIVITY DIAGRAMS AS FINITE STATE PROCESSES

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## HISTORY

- The INSPIRE project IST-10387-1999
- Supporting multiple views of business process models with a single core representation
- Views: RAD/AAD, IDEF0, IDEF3...
- Research on using process algebras for this purpose

## BUSINESS PROCESSES

- Business process:
  - a set of partially ordered activities carried out inside an organization, intended to reach a business goal and to produce a result of value for the customer.
- Notations for business processes:
  - High-level visual notations addressed to the business management community.
  - Low-level foundational notations addressed to the computer science community.

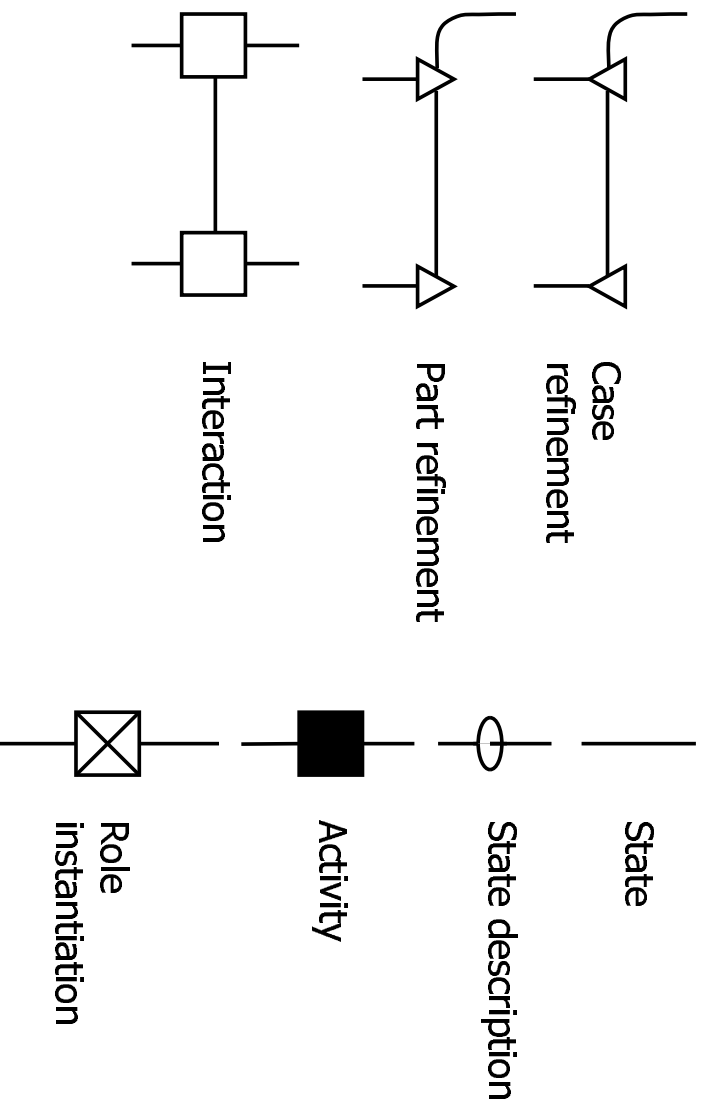
## MOTIVATION

- RAD is a popular high-level visual notation for capturing the dynamics and role structure of an organization.
- RAD has been used for modeling, simulation and enactment.
- There are no references in the literature reporting on the formal semantics of RAD.
- We are trying to bridge the gap between low-level and high-level formalisms for business process modeling.

## RAD NOTATION (I)

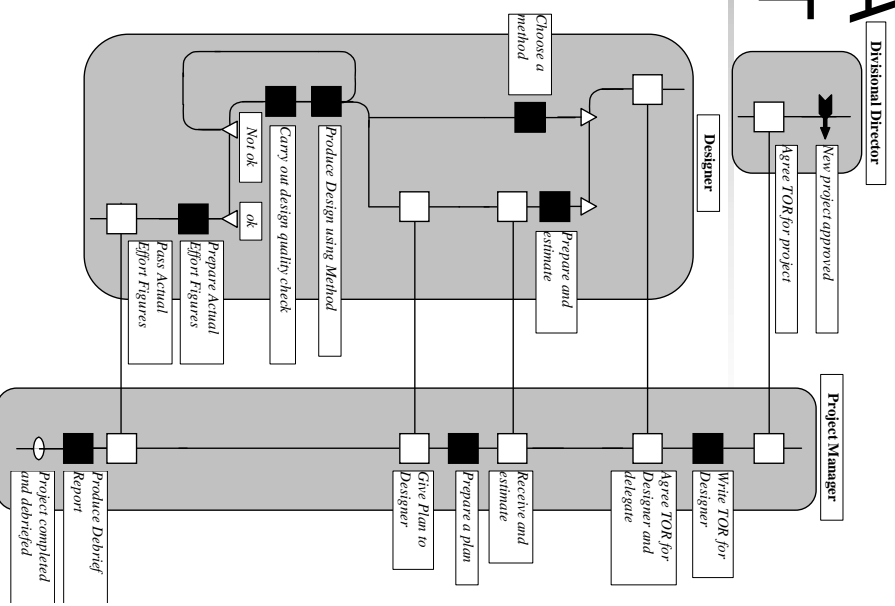
- *Roles* group together activities into units of responsibility. A role has one or more execution threads containing *case refinements* and *part refinements*.
- *Activities* are the basic blocks of a role. They can be carried out in isolation or may require coordination with activities in other roles. In this case they are called *interactions*.
- *External events* are points at which state changes occurring in the process environment influence on our process.

## RAD NOTATION (II)



# CARRYING OUT A DESIGN PROJECT

- Three roles:
  - Divisional Director
  - Project Manager
  - Designer



# FSP PROCESS ALGEBRA

- FSP is an algebraic specification technique of finite state labeled transition systems.
- The tool Labeled Transition System Analyzer – LTSA that allows:
  - Writing FSP specifications
  - Generating the underlying LTS
  - Verification of safety and progress properties

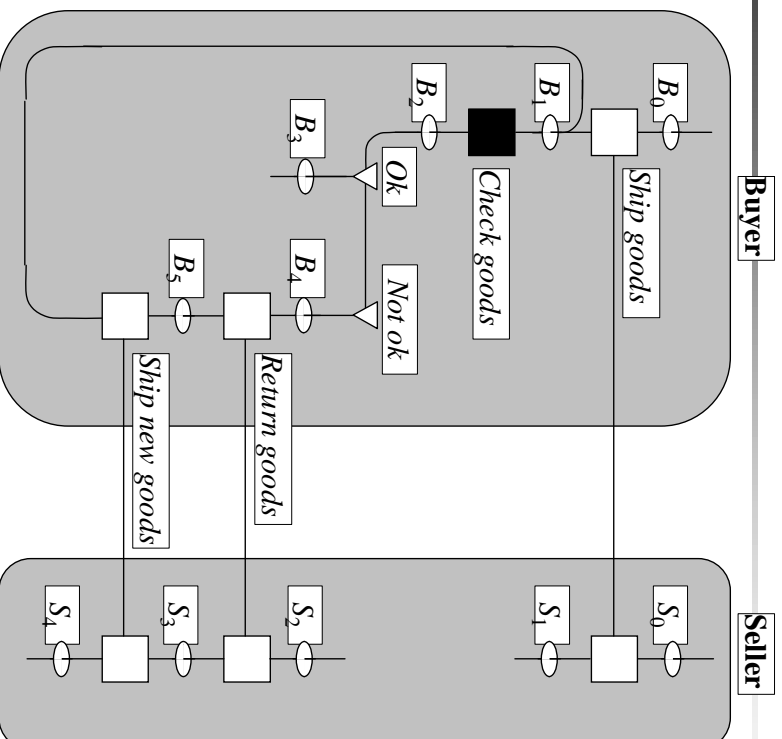
## FSP OPERATIONS

- Notation:  $a$  is an action,  $P$  and  $Q$  are process terms,  $H$  is a set of actions and  $L$  is a re-labeling function
- Operations:
  - Prefix:  $a \rightarrow P$
  - Choice:  $P \mid Q$
  - Parallel composition:  $P \parallel Q$
  - Hiding:  $P \setminus H$
  - Re-labeling:  $P / L$
  - Definition:  $A = P_A$

## FSP PROCESSES

- Primitive (sequential) processes
  - Is defined using a set of local definitions in which the process term is either END or choice.
- Composite processes
  - Is defined with a definition in which the process term is a parallel composition, a hiding or a re-labeling.
- Maintaining this distinction assures that the resulting models are finite.

# INTERACTION BETWEEN A BUYER AND A SELLER



## MODELLING RAD STATES

- State lines:
  - $NL = L0,$
  - $L0 \equiv (i \rightarrow L1|end \rightarrow ENND),$
  - $L1 \equiv (s \rightarrow o \rightarrow L0).$
  - $||L = NL \setminus \{s\}.$
- End states:
  - $E = (i \rightarrow E|end \rightarrow ENND).$
- Single instance start states:
  - $SS = (o \rightarrow end \rightarrow ENND).$
- Multiple instance start states
  - $SM = (o \rightarrow SM|end \rightarrow ENND).$

## MODELLING THE BUYER

$$\begin{aligned}
 & \| B_0 = SS/\{sg/o\}. \\
 & \| B_1 = L/\{\{sg, sng\}/i, cg/o\}. \\
 & \| B_2 = L/\{cg/i, \{n\_ok, ok\}/o\}. \\
 & \| B_3 = E/\{ok/i\}. \\
 & \| B_4 = L/\{n\_ok/i, rg/o\}. \\
 & \| B_5 = L/\{rg/i, sng/o\}. \\
 & \| Buyer = (B_0 \parallel B_1 \parallel B_2 \parallel B_3 \parallel B_4 \parallel B_5).
 \end{aligned}$$

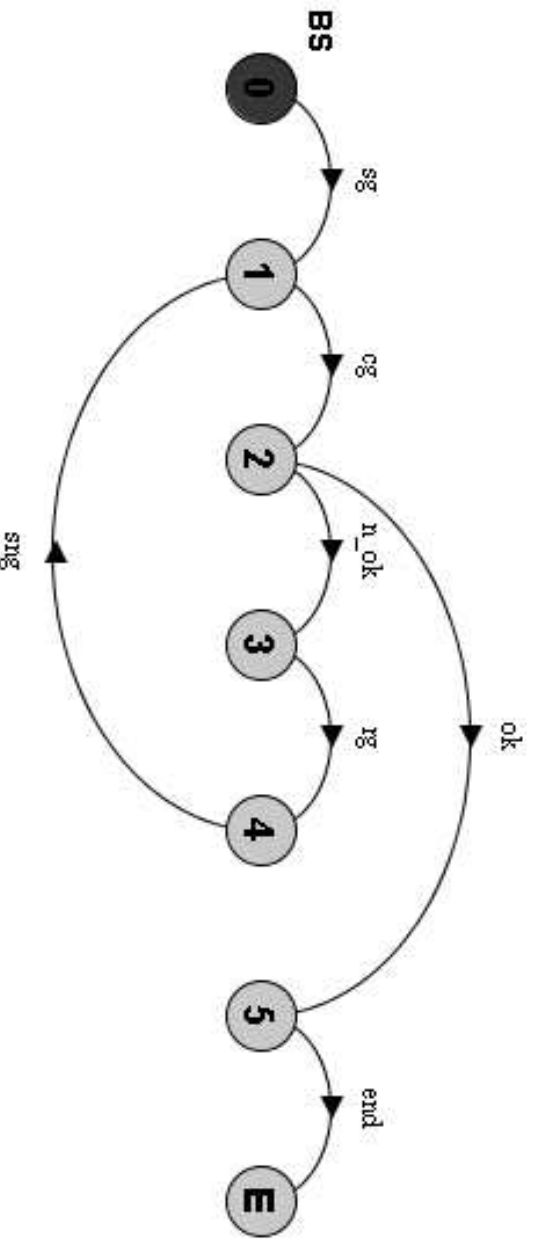
## MODELLING THE SELLER

$$\begin{aligned}
 & \| S_0 = SS/\{sg/o\}. \\
 & \| S_1 = E/\{sg/i\}. \\
 & \| S_2 = SM/\{rg/o\}. \\
 & \| S_3 = L/\{rg/i, sng/o\}. \\
 & \| S_4 = E/\{sng/i\}. \\
 & \| Seller = (S_0 \parallel S_1 \parallel S_2 \parallel S_3 \parallel S_4).
 \end{aligned}$$

# MODELLING THE INTERACTION BUYER-SELLER

$\|System = (Buyer \parallel Seller).$

# LTS OF THE BUYER-SELLER EXAMPLE



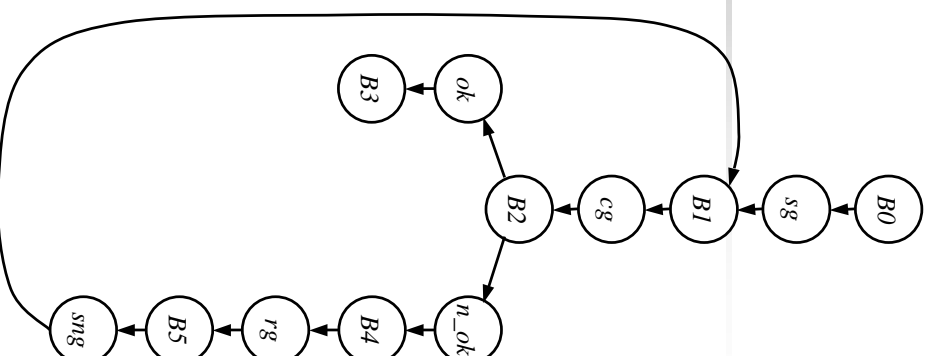


# RAD MODELS AS GRAPHS

- *RAD model* = a set of role models
- *Role model* = a bipartite directed graph with nodes partitioned into states and actions.
- States = states in a role
- Actions = activities, external events, case refinements and part refinements
- A case refinement with  $n$  alternatives  $\Rightarrow n$  actions
- A part refinement  $\Rightarrow$  at least one action node. If two or more threads originating from the same part refinement recombine together  $\Rightarrow$  a new action is added.

## EXAMPLE

- The graph for the buyer in the buyer-seller interaction example
- Uppercase nodes are *states* and lowercase nodes are *actions*.



## CLASSIFYING STATES

- It is needed for the mapping algorithm
- States are partitioned into:
  - Labeled state lines
  - Unlabeled state lines
- Start states = states with no in-coming edges
- End states = states with no out-coming edges

## MAPPING STATES

- An unlabeled state line  $s \Rightarrow$   
 $P(s) = L/\{E^{-1}(s)/i,E(s)/o\}$
- A labeled state line  $s$  with label  $l \Rightarrow$   
 $P(s) = NL / \{E^{-1}(s)/i,l/s,E(s)/o\}$
- A single instance start state  $s \Rightarrow$   
 $P(s) = SS/\{E(s)/o\}$
- A multiple instance start state  $s \Rightarrow$   
 $P(s) = SM/\{E(s)/o\}$
- An end state  $s \Rightarrow$   
 $P(s) = E/\{E^{-1}(s)/i\}$

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input a RAD model  $\mathcal{R}$ 
output a FSP model  $\Delta$  for  $\mathcal{R}$ )
 $\Delta \leftarrow$  definitions of processes  $L, NI, E, SS, SM$ 
 $PN \leftarrow$  a new composite process name
 $F_1 \leftarrow PN+' = ('$ 
for each  $R \in \mathcal{R}$  do
     $RN \leftarrow$  a new composite process name
     $F_2 \leftarrow RN+' = ('$ 
    for each  $s \in S(R)$  do
         $(SN, F_3) \leftarrow SFSP(s)$ 
         $\Delta \leftarrow \Delta \cup \{F_3\}$ 
         $F_2 \leftarrow F_2 \cup \{SN\}$ 
        if  $s$  is not the last in  $S(R)$  then  $F_2 \leftarrow F_2+' ||'$ 
         $F_2 \leftarrow F_2+' )'$ 
         $\Delta \leftarrow \Delta \cup \{F_2\}$ 
         $F_1 \leftarrow F_1 + RN$ 
        if  $R$  is not the last in  $\mathcal{R}$  then  $F_1 \leftarrow F_1+' ||'$ 
         $F_1 \leftarrow F_1+' )'$ 
     $\Delta \leftarrow \Delta \cup \{F_1\}$ 

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## POSSIBLE USES OF THE FSP MODEL

- Interactive simulation:
  - Trace into a process by executing it step by step.
- Static verification:
  - Check if the model has some desired correctness properties. This can be done using model-checking. LTSA has the possibility to check *safety* and *progress* properties.
- Dynamic simulation:
  - Estimate performance indicators of a business process: degree of utilization of resources, cycle times, waiting times, throughput, a.o.

## RELATED WORK

- Mapping RAD to DES (Martinez-Garcia, Warboys, 1998)
- Modeling PIF using CCS (Schroeder, 1999)
- Modeling workflow processes using FSP (Karamanolis, Giannakapoulou, Magee, Wheeler, 1999)
- Modeling UML AD with FSP (Rodrigues, 2000)
- Modeling UML AD with CSP (Bolton, Davies, 2000)
- Modeling RAD using PEPA (Badica et Fox 2001)

## CONCLUSIONS

- RAD notation has a natural mapping to FSP.
- The mapping has been presented as a translation algorithm.
- The resulting FSP model can be used to derive a LTS computational model of a business process.
- The computational model is useful for checking the behavior of the business process before having actually to implement it in practice.